IntegratedMobilityConceptsforWirelessLANs

MatthiasSiebert ¹,ArndtKadelka ¹,MatthiasLott ²,MartinWeckerle ²,SvenHischke ³

¹AachenUniversity,ChairofCommunicationNetworks,{Ma tthias.Siebert|Arndt.Kadelka}@comnets.rwth-aachen.de ²SiemensAG,Mü nchen,Germany,{Ma tthias.Lott|Martin.Weckerle}@icn.siemens.de ³T-SystemsNovaGmbH,Darmstadt,Germany,Sven.Hischke@t -systems.com

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Abstract

Thefocusofthispaperisontheintegr atedIPmobilitysu pportwithinmicrocellularsystemsasanextensiontoglobal mobilityenvironments.Sincemacromobility,mainlyrea 1izedwiththehelpofMobileIP, is not able to support highly dynamicscenarios, novel concepts have to be deve loped. Differentsolutionsareproposed within this paper, e.g. for fasthandoverinWLANs, which are co nnectedtoanIP basedinfrastructure.Forexpeditedhan doverexecutiona newconceptforbackbonesignalingisproposed.Ite xplores theinterfacebetweenthe linkandnetworklayer,whichwill bedescribed,too,andwhichhasbeenspecifiedindetailfor HIPERLAN/2.Withthea pproachesdescribedinthispaperit canbeexpectedthatHOcanbesupportedinanefficientway notsacrificingtherareradioresourc e.

1.INTRODUCTION

Inordertosatisfythemobileuser'sdemandforhighdata rateswithQualityofService(QoS)support,innovativeco nceptsneedtobeintroduced.Current3Gcellularsystems onlypartlycovertherequiredneeds. Therefore, one promi singapproach, which is also followed within the Information SocietyTechnologies(IST)projectMIND(MobileIPbased NetworkDevelopments) [1]aimsati ntegratingWireless LANtechnologyasanaccessnetwork, complementary to 3G cellularsystemst ogetherwiththeirwired, mainlyIP -based fixedbackboneconnections.Nexttohighdatarates,which areprovidedbyWLANs,fastandseamlesshandover(HO) supportisofessentialinterestifusersmovewhileaconne ctionisestablished.Th isb ecomesagreatchallengesincethe coverageofWLANsisrestrictedduetothehighfrequencies theyoperateat(us uallyatseveralGHz).

CurrentresearchactivitieswithinEuropeanprojectsandr espectiveworkinggroupsoftheIETFfocusonenhanced mobilitysupport(e.g.WINEGLASS,MobyDick,IETFm obileIP,...).Sincemostoftheexistingaccessnetworksalready havetheirownsupporttohandlemobilitybetweenAccess Points(APs),projectslikeWINEGLASS [2]employana rchitecturethatreliesontheexistingmobilitymechanisms providedbythea ccessnetworks toprovidemicro -mobility. OtherprojectslikeMobyDick [3]ha ndlemobili tyatIP -level orabove [4].

However,noneoftheapproachesincorporatesdifferentm icromobilitysupportschemesinordertointegratetheminto anoverallconcept.Bymeansofcombiningtheadva ntageous aspectsofeachpropos edconcept,amoresophist icatedand efficientHOsupportcanbeachieved.E.g.informationgat heredwhileapplyingoneHOschemecanbeusedtospeedup anotherHOscheme.Insuchawayaco mplementaryoverall conceptisachieved,whichwillbepr esented inthispaper.

StructureofthePaper

TodescribeourapproachesfortheHOHIPERLAN/2(H/2) [5][6][7]isusedasexa mplesystem.Itprovideshighdata rates, meansforQ oSsupportanditis assumed as compl ementarywirelessaccesstechnologyofUMTSby ETSI/BRANand3GPP [8].Afterintroducingthemobility controlofH/2inSection2,abriefoverviewisprovidedof theinterfacebetweenthel inklayerandnetworklayerinSe ction3.ThisinterfacehasbeenspecifiedindetailforH/2and providesallrelevantprimitivestosupportefficientHOco ncepts.InSection4twoconceptsarepr esentedhowtheHO andthusthemobilitysupportcanberea lized.Themainidea istoshifttheinform ationtransferofconnectionrelateddata, whichneedstobepresentatthenewAP,fromtheairinte rfacetothewiredbackbone.However,aminimumamountof signalingviatheairinterfacealwaysneedstobedon e.Co nsequently.thedifferencebetweenthetwohereindescribed HOstrat egiesliesintheamountofdatathatneedstobe transmittedviatheairinterface.Theadvantagesofeacha pproacharepointedoutandawayofcombiningbothmethods ino rdertoac hievebestpossibleHOsupportispresented.In Section5wedescribeawayofefficientlybuildingupad atabase, which is necessary to support fast HO. This database alsodemonstrates the complement character of our co ncepts, sinceitissetupwiththe helpoftheoneHOtypeandused bytheotherone.InSe ction6wepresentsimulationresults that will show the ffectiveness of our proposal. As already mentionedabove,H/2ischosenasWLANreferencesta ndardfordemonstratingthekeypointsofthecon cept.Ho wever, the described method is not limited to H/2 and can be appliedtomobilitysupportwithinWLANstandardsinge neral.

2.MOBILITYCONTROL

ThecurrentH/2specificationdealswiththreetypesof HO:sector,radioandnetworkHO.Whilethefirs ttwoones canbecompletelyservedwithinoneH/2accesspoint,the latteronerequiresthesupportoftheaccessne twork.The proceduresforthesectorandtheradioHOarespecifiedin [7].Withinthispaperonlyashortove rviewwillbegiven. Themainobjectiveistodescribetheproc edures, which are necessarytoprovideanetworkHO.SincetheMobileTe rminal(MT)leavestheservingareaofaRadioLinkControl (RLC)instance, ane tworkHOinvolvesalsohigherlayers, whichareco nnectedtotheH/2DataLinkLayer(DLC)byan appropriate convergence layer (CL). This CL is important for theinterconnectionofaspecifichigherlayerpr otocollike e.g.Ethernet,IEEE1394orIPwiththeDLC.Thetaskofthe CLtherebyisto adapttheservicerequirementsoftherespe ctivehigherlayertobearerservicesofferedbyH/2.Thisi ncludesboth,logicalQoSmappingaswellasadjustmentof potentialvar iablehigherlayerpacketsizestothefixed packetlengthsofH/2.Inthenotati onoftheCL,aprefixis addeddependingonthehigherlayertheCLisconnectedto: E.g.tomaintainassociationandconnectionparameters.sp ecificsignalingviatheIPbackbonemaybeneeded.Thisr equiresinterworkingwiththeIPlayerbymeansofan IPco nvergencelayer(IP -CL).Amoredetaileddescriptionofthe IP-CLwillbegiveninSe ction3.

3.INTERLAYERCOMMUNICATION

ClassicalWLANsystemsdonotspecifyhigherlayer functionalities.Sincetheirfocusistoserveasradioaccess network,th osefunctionalities,e.g.routingoraddressresol ution,needtobeovertakenbyarespectiveIntelligentNe twork(IN),whichisalsoconnectedtoorpartofafixed backbonearchitecture.Inordertoproviderequiredse rvices totheupperlayers,theinfo rmationtransferbetweenthe lowerlayersoftheWLANsystemandthehigherlayerofthe INneedstobeorganized.

WithintheISTprojectBRAIN [9],agenericinterfaceIP ₂W [10]wasspecifiedthatisab letoovertakethistask,seeFi g-ure1.

ThegenericdesignoftheIP ₂Winterfaceallowstheinte rworkingofvariousWLANswithrespectiveINs.Inorderto facilitatethisinterworking,therespectiveWLANsy stem onlyneedstoincorporateaspecificCLtha tisabletounde rstandtheIP ₂Wd irectives.

TheCLitselfisdividedintoaCommonPartConvergenceSublayer(CPCS)andaServiceSpecificConvergenceSublayer(SSCS).ThecurrentH/2specificationsalreadyprovideaCPCS[18]aswellasIEEE1394andEthernetSSCS[19][20]whereasBRAIN/MINDfocusesonthe

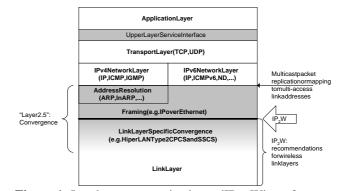


Figure 1: InterlayercommunicationandIP 2Winte rface

specificationofaninnovativeSSCSded icatedtothedirect supportofIPtrafficinamobileenvironment [21].Thiswill allowprovidingtherequestedIP ₂Wservicetothenetwork layeroneachsideoftheairi nterface.

Withintheproposedconceptsformobilitysupport,theH/2 IP-CLwillbeemployedandslightlyenhancedinorderto connectH/2totheINandtoenableinformationexchange betweenoldandnewAPincaseofaHO.ThoughH/2is chosenasanexemplaryWLAN,theconceptsarenotfocused onthisstandard, sincetheemploymentofthegenericIP₂W interfaceallowsama ppingtootherWLANstandards,e.g. IEEE802.11,too.

4.INTERWORKINGWITHIPMOBILITY

WLANstandardslikeH/2shallsupportenddevices' mobilitywithinandoutsideofactivecommunic ationp hases. AterminalthatisassociatedtoanAPresidesinthea ctive phase,regardlesswhetheritistransmittinguserdataornot. Duringtheassociationprocedure,theMTregisterswithinthe localaccessnetworkwiththehelpoftheAP.Inthefirst case,whentheMTregisterswithinaforeignne tworkandhas noactivecommunicationyet,thehomene tworkneedstobe informed.Forthispu rpose,theMobileIP -protocol [13]can beused.

Thefollowingappro achintegratesH/2 -andIP -mobility.We thereforeselectahierarchicalmobilitymodelanddi stinguish betweenglobalandlocalmobility.Withinthisnetwork model,aMobileIPcapablenetworkprovideswide(global) areamobility,whereaslocalwirelessac cessnetworksasco n-sideredinMIND,handlelocalmobility.Withinthelocal area,theMobileIPprotocolisnotcap abletobeusedsinceit isoptimizedforslowlymovingMTsandbecomesinefficient inthecaseoffrequentmigrations [12].Forthatreason,more recentapproacheshavetriedtotakeintoaccountthelocal characterofmobilitytoavoidtime -consumingre -routing

overthehomeagent.ExamplesareCe llularIP [14], HAWAII [15]orHierarchicalMobileIP [16].

TheefficientsupportoflocalmobilitybetweendifferentAPs demandsforacloseintegrationofH/2 -mobilitypr otocols and the IP -protocol. Differenta pproaches for local mobility protocolsassumethetransmissionofroute -updatepackets viaanactiveradioconne ctionbetweentheMTandtheAP. Therefore, in the case of an H/2HO firstly the DLC proceduresforassociationandconnectionestablis hment needto berunthrough. Thus, anupdate of the rou tingpath withintheaccessnetworkcanonlybedoneafte rwards.This time-consuming, sequential runoff can be avoided, if the newAPtriggerstheIPmobi lityprocedureassoonastheMT hasregistered.Userdata tobetransmi ttedtotheMTwillbe bufferedbythenewAPuntiltheDLC -connectionsareesta blished.

Anotherimportantadvantageofthissolutionthatdisti nguishesthisapproachfrompreviousproposalsis,thatthe protocolsforlocalmobilitysupport neednottobeimpl ementedintheMTs.Thus,thedifferentlocalnetworksmay evenemploydifferentmobilityprotocolswhereasthei mplementationwithintheMTsisnotaffected. Thiscloseintegrationofthemobilityprotocolsisthetaskof theIP -CL.The networkHOprocedureoftheRLC -protocol thereforeconveysdifferentpossibilities,depen dingonthe supportedserviceswithintheaccessne tworks.

Thefollowingsectiondescribestwopossiblevariants:Ne t-workHOwithandwithoutsignalingoverthebac kbonene t-work.

4.1Networkhandoverwithoutsignalingoverthebac kbone

IfaMTinitiatesaHOwhileanactiveconnectionise stablished,theinformationaboutthecurrenttransmi ssion statusneedstobereportedtothenewAP.Iftheaccessne tworkdoe snotsupportthetransferoftheinformation,they needtobere -negotiatedbetweenMTandthenewAP.Fi gure2showstheschematicsequenceofthenetworkHO. Sincetheconnectionrelateddataistransmittedviatheairi nterface,therespectivesub -proceduresforcallestablis hment needtoberunthrough.

Thefollowingstepsareexecutedinchronologicalorder:

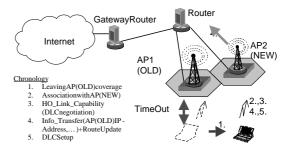


Figure2: Handoverwithoutsignalingoverthebackbone

- 1. LeavingofthecoverageoftheoldAP(AP1(OLD)).
- AssociationoftheMTatthetargetAP(AP2(NEW)).Thereby theMTtransmitsitsoldMAC_ID(MAC_ID1),whichwas usedintheongoingconnectionwithAP1(OLD).Additio nally,theDLCaddressofAP1(OLD),theso -calledAccess PointIde ntifier(AP -ID1)istransmitted, too.Asaresponse, AP2(NEW)assignsanew,localMAC_ID(MAC_ID2)tothe MT.
- Afterwards,theHO_Link_Capability -sub-procedureiscalled tonegotiatetheDLCconfigurationparametersb etweenMT andAP2(NEW).Besidesthis,AP2(NEW)d etermineswhich othersub -proceduresmightneedtobecalled.
- 4. TheInfo_Transfer -sub-procedureisusedbytheMTtoco nvev information about its own addresses, like the Logical LinkAddress(LLA)e.g.inEUI -64-formatandtheIP -address. Whererequired, the AP1 (OLD) IP -addresscanalsobeco nveyed.AP2(NEW)hereafterconfirmstheinformationtran sfer withitsownIP -address.Havingallrelevantinform ation,the IP-CLofAP2(NEW)willconveythisinform ationtotheIP layer,whichthencaninitiatetheroute -updatepacketof theIP mobilityprotocol.Such,theinternalroutingtablesofthe routerareu pdated.
- Almostinparalleltotheroute -updatesignal,theDLCconne ctionsviatheairinterfaceareestablishedafterwards.TheMT's IP-CLreportsthenewconne ctionstoitsIP layer.
- 6. TerminationoftheHOprocedure.

Theabovedescribedproceduredoesnotexplicitlyinform AP1(OLD)abouth eongoingHO.Therefore,itkeepsthe associationwiththeMTuntilatimeoutcomesupandforces theAP1(OLD)tostartthedisassociationprocess.Userdata, whicharrivesattheAP1(OLD)withinthisp eriodcannotbe forwardedandgetslost.

Thisisal soindicated within Figure 3, which illustrates the involvement of the IP -CL within this HO type: Theold AP is not informed about the change of association of the MT and thus does not know, whether the MT has left its cove rage area, wasswitched off or su ffers from an insufficient radio link quality. Therefore, it will receive a time outsi gnal and release resources, which we reallocated for the MT's conne ction.

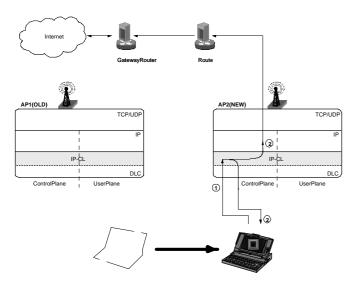


Figure3: IP-CLwithinhandoverprocess:Nosignalingover thebackbo ne

Sincethenewpointofattachmentisnotknown,notunne ling ofmisrouteddatapacketsispossible.WithinthisHO,only theIP -CLsoftheMTandthenewAParei nvolved.The aforementionedactionsareperformedwithintheco ntrol planeoftheMTandA P2(NEW).Thesubsequentdatatran sferviatheuserplaneisnotincludedinthefi gures.

TheIP -CLwithinthistypeofHOmainlyrestrictsonfo rwardingthereceivedinformationtotheupperlayers.Thus, noadditionalchallengesneedtobefulfilled.

4.2Networkhandoverwithsignalingoverthebac kbone

There -negotiationoftheconnectionrelatedparam eters viatheairinterfaceduringaHOresultsinanappreciablei nterruptionofdatatransmission.Thus,inthefo llowing.itis assumed that this in formation is transferred to the target AP (AP2(NEW))viathefixedwirelineaccessnetwork,cf.Fi gure4.Forthispurpose,theintegrationofacontrolprot ocol, located within the IP - CL of each AP, is necessary. We therebyassumethatdatatransfervia thefixednetworkis muchfasterandmorerel iablethanviatheairinterface.In thefollowingwethereforealsor efertothetwoHOtypesas "slow" and "fast" HO, since it is very likely that the HO type withsignalingoverthene tworkisfaster, becaus ethedatato betransmittedislessprobablycorrupted(noi nterferences). neitherthetimeconsumingaccessproceduresneedtoberun through.SimulationresultsinSection6supportthisstat ement.

TheproposedInter -APprotocolonlyconsistsofafew messagesandemploysastandardtransportprotocol, suchase.g. UDP [11]. The advantage of UDP incontrast to TCP is a shorter signaling period, but no error protection is provided.

Figure4showstheshortenedHOprocedureatt heairinte rfaceandthecommunicationofthetwoAPsviathefixed wirelineaccessnetwork.Thefollowingstepsareexecutedin chronologicalorder:

- 1. LeavingofthecoverageoftheoldAP(AP1(OLD)).
- 2. DuringtheHOAssociationtheMTprovidesitsorig inal MAC_ID1andtheDLCaddressofAP1(OLD)(AP -ID1)to thenewAP2(NEW).ItthereafterreceivesitsnewMAC_ID2 fromthenewAP2(NEW).Inthefollowing HO_Link_Capabilitysub -procedure,AP2(NEW)indicatesthe shortenedHOproc edure.
- 3. BasedontheAP -ID1, AP2(NEW)determinestheIP -address ofAP1(OLD).ForthispurposeitisassumedthatAP2(NEW) holdsalookuptable, within which incorp oratedAP -IDsare mappedtotheirrespectiveIP -address.Thislookuptablecan either beman a ged within the IP-CLorwi thintheIP -layer. Generation and update of this table is the task of address res0lutionprotocols.AnexampleisdescribedinSection5.Afte rwards, AP2(NEW) informs AP1(OLD) about the ongoing HObysen dingamessageviathefixednetwork, usinga higherlayertransportprot ocol(e.g.UDP).Themessage shouldcontaintheparametersAP -ID1andMAC_ID1aswell astheAP -ID2andtheAP2IP -address.
- AP1(OLD)replieswithamessage,whichbesidestheMTa ddresses(EUI -64andIP)includesinformationabouth eexis tingDLCUserConnections(DUCs).Thedescribingparam etersDUC -descr,e.g.DLCC -IDsorCL -CONN-ATR(datar elatedtotheCLwhichistransmittedtransparentlytoRLC)are identicaltothoseoftheRLCprotocol[7].Userdata,stillmi sroutedfromth eroutertoAP1(OLD)canthenbeforwardedto thenewAP2(NEW)bymeansofIPtunneling.
- 5. OnhavingreceivedallthenecessaryDUC -descr,thesetupof theDLCconnectioncanbeco mpleted.
- 6. TheIP -CLofAP2(NEW)subsequentlyinformsaboutthe newconnectio ns.AP2(NEW)generateswiththehelpofthe received information from AP1(OLD) aroute -updatepacket (not shown in Figure 4) and reports the new point of attac hment of the MT to the router.
- 7. TerminationoftheHOprocedure.

Figure5illustratesthetask oftheIP -CLwithinthisHOtype. Itcanbeseen,thattheAP2(NEW)nowretrievesinform ationrelatedtotheongoingconnectionoftheMTviathewire lineconnection.

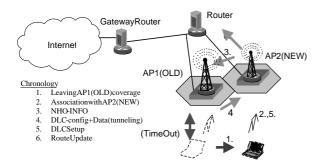


Figure4: Handoverwithsignalingoverthebackbone

Therefore, it firstly needs informat ion about the old AP (AP-IDold)andthemobileidentifieratthisoldAP (MAC ID). This minimum amount of informational ways needs to be transmitted by the MT via the air interface during theHO_Link_Capabilitysub -procedure.Havingthisinfo rmation, AP2 (NEW) now is able to contact AP1 (OLD). For thispurposeasecond, independent instance (HO -INFOi nstance, AP2b) is initiated within AP2 (NEW). Insuchaway, allrelevantinformationthatispr esentinoneoftheinstances ofAP2(NEW)maybeaccessedbyt heotheri nstance,too. ThedistinctionAP2a/bhereinisintroducedasalogicalsep arationtoindicatethattherearetwoprotocolsrunning(which arehandledbytwoinstancesatAP2(NEW)),oneofthem servingthecommunicationbetweenAP2(NEW) -MTandth e otheroneservingthecommunicationbetweenAP2(NEW) AP1(OLD).

TheIP -CLoftheHO -INFOinstance(AP2b)reportsther eceived information to the upper layers and demands for pr0visionofthemissinginformationabouttheoldlayer2co nnectionfrom the old AP. Hereupon, AP2 bestablishes a higherlayerconnection(e.g.UDP)withAP1(OLD)and transmitstheparameterstransparentlytoAP1(OLD).InAP1 (OLD)therequired information is gathered from the IP -CL (see4,5inFigure5)andsentbacktoAP2b.A P2bfinally providestherequestedinformationtotheIP -CLwhereitcan beaccessedbyitscallinginstance(AP2a).AP2athenisable tocontinuewiththeDLCsetupproc edure(8).Atthesame time,theAP2aIP -CLinformstheupperIPlayeraboutthe newlyarrivedandassociatedMT.TheIPlayerthensends theroute -updateinformationtotherouter.Adetailedd escriptionoftheprimitives, which are used within the info rmationexchangeprocedure,t ogetherwiththeirattributescan befoundin [10] and [17].

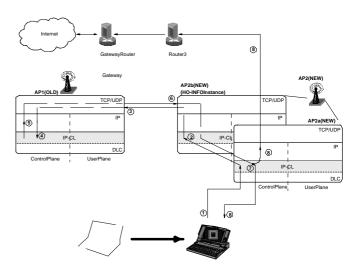


Figure5: IP-CLwithinthehandoverprocess:Signa lingover thebackbone

TheIP -CLwithinthistypeofHOdiffersinitsco mplexity comparedtotheoneintheprevioussection .Forthewireline supported signaling, the IP -CL needs to fulfill add itional taskslikemanagementandupdateofthedat abasepresented inthenextsection.Onefurtherdiffe renceisthepointoftime atwhichtheIP -CLinformstheIPaboutthenewconne ction (andthustriggerstheIPlayertosendtheroute -updatepacket totherouter).InthefirstHOtype,thisisdoneassoonas possible, which means even before the complete DLC co nnectionisestablished.Inotherwords:HOsonlayer2and layer3are executedinpara llel.Thisisnecessarytoavoidthe lossofmi sroutedIP -packets.InthesecondtypeofHO,the IP-CLfirstlywaitsuntilthecompleteDLCconnectionise stablishedandthenreportsthenewconnectiontotheIP -layer. Thelayer2HOhasc ompletelyfinishedbeforethelayer3 HOisexecuted.Itisnotnecessarytoexecutetheminpara 1lel, sincemisrouted packets are end by means of tunne 1ingandstoredinAP2, thus no packets getlost.

5.DYNAMICCOMPOSITI ONOFATOPOLOGY DATABASE

Theaboveintroducedprocedureforapresumedfaster HOexecutionwithsignalingviathefixedne tworkneedsto beawareoftheIPaddressesofneighboringAPstoa pplythe Inter-APprotocol.EachAPadministersasimplifiedtopo 1_ ogydatabasewiththeIPad dressesofitsHOpartners,the precedentservingAPofaMT.Thisdatabaseisstructured likealookuptable,whereinama ppingfromIP -addressesto DLC-addressesandviceversaispe rformed.Thedatabase caneitherbemanagedwithintheIP -CLorintheI Player. Oneargument for the latter surely is, that IPa ddressesare layer3specificandthusshouldbeadministeredthere.Ho wever, it is proposed to ocate it in the IP -CL, since this means that the decision about the type of HO to bech osen does not need an intervention of layer 3 anymore. Add it ionally, it was stated that AP2 aand AP2 behare the same IP -CL space. Since the information about the IP address of AP1 (OLD) needs to be transferred to the IP layer of AP2 b, it is favorable if the data base is in the common ly shared IP -CL. The following method describes away to dynamically setup and update such adata base:

- 1. TheMTregistersatthenewAPwhileconveyingtheAP -IDof theprecedentAP.
- 2. ThetargetAPdecideswhichkindofHOprocedureisa pplied. thismeanseitherthe"slow"onewithoutfixednetworksu pport,orthepresumedfasteronewithfixednetworksu pport.If thenewAPisnotabletoperformamappingb etweentheAP -IDandtherespectiveIPaddress, it will decide to carry out the slowHO execution,whichmeansalldatai ncludingtheIP addressoftheprecedentAPisre -transmittedviatheairinte rface,thustheinternaldat abasecanbeupdated(seeright branchinFigure6).
- 3. If the AP -ID is in the database, the AP may decide to use the HOwith signaling over the wireline access network. Ho we ever, if the present database needs to be updated, the AP can decide to use the other type of HO neve rtheless.
- WithintheslowHOprocedure,theAPreceivesinform ationto update/setupitspersonaltop ologydatabaseandstoresther espectiveco upleofAP -IDandIP -addresstherein.

If MTs from the same precedent AP enter the cove rage area of an AP, all the relevant address information is available and the "fast" HO procedure can be executed. For the pu pose of updating the database, the slow procedure can be demanded periodically, see Figure 6.

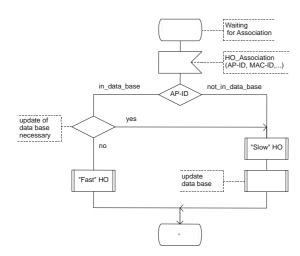


Figure6: Dynamiccompositionofatopologydatabase

6.ANALYSIS

Inthefollowingsectionweinvestigatetheinfluenceof thedifferentHOprocedureson thepacketdelay.Theafor ementionedconceptshavebeenmodeledand implementedandwillbeevaluatedbymeansofcomputer simulations.

ScenarioDescription

WeassumeasimplescenariowithoneMTmovingfromone AP(AP1)toanotherAP(AP2),asshownw ithinFigure2 andFigure4.Sinceourfocusisbasedontheprotocolcaused delaytimes, we presume a perfect link without any disruption orlossofpacketsontheairinterface.Datapacketsonlyget lostincaseofbeingmisroutedduringtheHOprocess andits subsequentassociationprocedure.SincenoAutomaticR epeatRequest(ARQ)isapplied,thosepacketswillnotbere transmitted.ForallsimulationsweconsidertheMTtomai ntainanactivebi -directionalco nnectionwithsymmetric DUCsloadedby64k bit/s,respe ctively2Mbit/s -Poisson sources.Onlypacketsthatreachthereceiverwillbeeval uated, thus packets that getlost do not contribute to the delay evaluation.

Handoverwithoutnetworksupport(slowHO)

Westartbyexaminingtheprotocolcaused delaytimes withinthenon -backbonesupportedscenario.AssaidinSe ction4.1,alldatarelatedtoDLCassociationandconnection needstobere -transmittedviatheairinterface. Figure7pointsupthemessagesandPDUsthatneedtobe



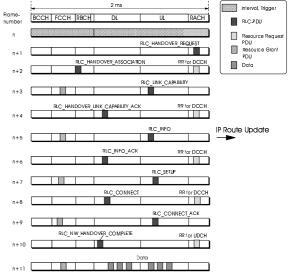
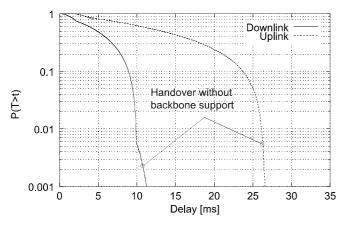
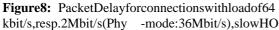


Figure7: RLC-messagesforHOwithoutbackbone support

Figure8showsthecomplementarydistributionfunction (CDF)ofthesimulateddelaytimesforthisHOtype, regardingup -anddownlinkseparately.





ForthesimulationsinFigure8wetookPoissonsourcesge neratingtrafficof64kbit/sresp.2Mbit/s.SincetheCDFsof thedelaytimesof bothtrafficloadsarethesame,onlytwo curvescanbeseenhere. Thereason for this congruence lies inthechosenPhy -mode16QAM3/4withitsnominaldata rateof36Mbit/sthatisevenabletoservetheloadof2 Mbit/swithoutcausingadd itionaldelay. ThediscrepancyinthedelaytimesforULandDLisdueto thelossofDLpacketsduringtheHO,thustheyarenotco nsidered within the evaluation. Both curves show an asym ptoticbehaviortowards11ms(Downlink)and27ms(Uplink). Thiscanbeexplained withthehelpofFigure7.Assaidin theprevioussection, DL packets getlost during the HOpr 0cedure.ThishappensaslongastheIProute -updatepacket (seeFigure2and7)hasnotbeensent.ConcerningtoFi gure 7,atleastanother5MACframes(5x2m s=10ms)are needed, before earliest user data may be transmitted again. Fortheuplink, whereas no packets getlost the packets have tobecachedbytheMTandcannotbesentbeforethewhole associationprocedure isrunthrough. As indicated by Figure 7, this last satleast 10 MAC frames (10 x 2 ms = 20 ms < 27 ms). Herewenoteanothereffectthatcausesextradelay:Dueto thelongqueueintheMT,allbuffereddatacannotbesent withinoneMACframe, but additional frames are needed. FortheDLthiseffect doesnotappear, since themi srouted packetsarenotbufferedbutlost. Figure9showsthenumberoflostPDUsforbothtraffic

loads.Forthe64kbit/sloadweregisterameanlossof3 PDUs,whereasthemeanlossforthe2Mbit/sloadis79 PDUs.

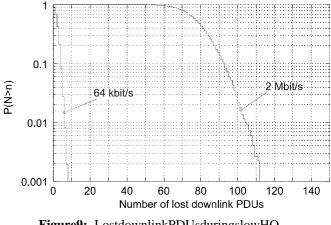


Figure9: LostdownlinkPDUsduringslowHO

Inordertohighlightthepreviouslydescribedeffectforthe ULdelaytimes, alowerPhy -mode, BPSK1/2, witha nominal datarateofonly6 Mbit/swaschosen, whereas the source with 2 Mbit/swaskept.

Figure 10 shows the result of this simulation. The delay for both connections increases, due to additional time needed to empty the buffer.

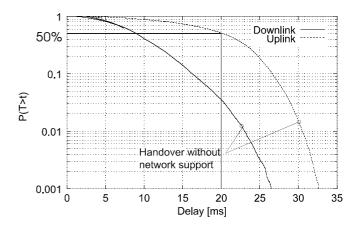


Figure10: PacketDelayforaconnectionwithloadof2 Mbit/s(Phy-mode:6Mbit/s),slowHO

Handoverwithnetworksupport(fastHO)

InthefollowingwefocusonthebackbonesupportedHO, presentedinSection4.2.Now,mostoftheassociationand connectionrelateddataistransmittedviathewiredbackbone from the oldAP1tothenewAP2,cp.Figures4and5.Only littleinformationneedstobeprovidedbytheMTandtran mittedviatheairinterface.

Figure11showsthemessagesandPDUsthatneedtob transferredforthisHO.

s-

e

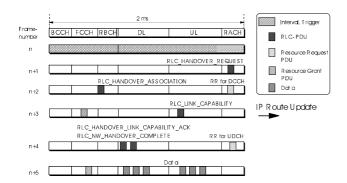
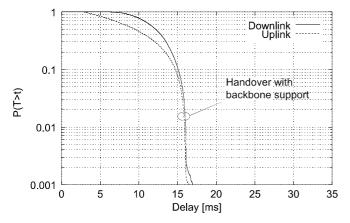


Figure11: RLC-messagesforHOwithbackbonesu pport, fastHO

ThemessageRLC_HANDOVER_REQUEST contains the DLC-ID of the old AP1 them oblew as connected to be fore. The new AP2 determines the respective IP - address with the help of the data based escribed in Se ction 5 and a cquire sall DLC related information from AP1. Comparing Figure 11 and Figure 7 we can see, that the HOnowise excuted si g-nificantly faster. Also the route - update message can be sent earlier.



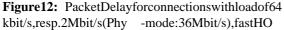


Figure 12 shows the simulation results for the same 64 kbit/s, 2 Mbit/straffic sources as for the previous simulations. Again, we apply first the highest Phy -mode 16 QAM3/4 with the nominal datarate of 36 Mbit/s. For the simulations we assumed that the data transfervia the wired back bone is su ficiently short and is executed in parallel to the DLC setup. As we can see, there is almost no discrepancy any more be tween the CDF sof UL and DL, since this type of HO allows tunneling of misrouted packets from AP1 to AP2, thus no (DL) packets get lost and now contribute to the evaluation. This is also there as on, why the total amount of time for DL packagesseamstodeteri orate(compareDLcurvesofFigure 8and12).However,comparingtheuplinkcurvesofboth Figuresdisclosestherealgain.

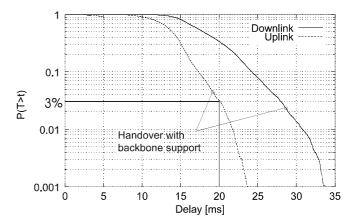


Figure13: PacketDelayforaconnectionwithloadof2 Mbit/s(Phy -mode:6Mbit/s),fastHO

Thesameisvalidifweregardalsotheheavilyloadedsc enario(load:2Mbit/s,Phy -mode:6Mbit/s)ofFigure13. Again,wecomparethe *uplink*curvesoftherespectivefi gures13and10,sinceth eyarenota ffectedbymi srouted packets. WhileforthefirstHOtypewithoutbackbonesu pporttheprobabilityforapackettosu fferfromadelayofe.g. morethan20msis50% (seeFigure10),therespectivepro babilityforapacketservedbytheotherHOtypeisonly3%.

7.CONCLUSION

nthispaper,aconceptforefficientandseamlessmov I ement between different network environments with a slittleinterventionfrom the user as possible is targeted. For that purposeagenericIP 2Winterfacehasbeendevelopedwithin theBRAINprojecttoc onnectdifferentkindoflinklayers andtheIPlayerandtoservetheIPlayerwithappr opriate triggersformobilitysupport. The applicability of this generic concepttoH/2resultsinanappropriateIP -CL,which supportstheHObetweendiffe rentAPcon trollers.This horizontalnetworkHOinvolvesi nteractionwiththeIP mobilityprotocolsandsignalingb etweenthelinklayerand IPlayerthroughtheIP -CL.Whereasformacro -mobility commona pproachescanbeapplied,forthemicro -mobility withintheacc essnetworkdifferentconceptsarediscussed. Especially, for fast HO procedures with r educedsignaling overtheairinterfaceanewapproachhasbeendescribedin thispaper.Adetailedd escriptionofthefunctionalityand proceduresinvolvedareexemplar ilypr ovidedforH/2. Toavoidtimeandresourceconsu mingauthenticationand registrationprocedureswhenamobilechangesthepointof attachment, an ovelex pedited signaling scheme over the

backbonenetworkisproposed.Bymeansofdiscoverypr 0cedures and caching each AP is a ware of the neighboring APsandallowstoexchangeuserandconnectionrel evant datafromtheoldtothenewAPinsteadofsi gnalingoverthe air.Aprerequisiteistheknowledgeofroutestothe neighboringAPs.Therespectiveprot ocolsneededtod ynamicallybuildupsuchatopologydatabaseineachAPco ntrollerhavebeendescribedinthisp aper,too.Bymeansof simulationswethereafterinvestigatedtheeffectsofbac kbonesignalingonthepacketdelay.Obviously,theaccele ratedHOprocedurewillbethebasisforapossiblesea mless HO.However.wehaveshownthatthefastHOisnota standalonesolution, since its accuracy depends on a dat abase, which is maintained up to date only due to the deplo ychawaybothconceptsco mentoftheslowHOtype.Insu mplementeachother.

Withthisnewconceptofsignalingoverthebackboneand buildinguptopologydatabasesaseamlessHObecomespo ssible,whichexploresthescareradioresourceinaneff icient wayatthesametime.

Sincetheb asicconceptbehindthenewproposedHOscheme forhorizontalHOasdescribedinthispaperisindependent oftheaccesstechnology,itisexpectedthatforthevertical HOthisconceptmayalsoplayafundamentalroleforeff cientandseamlessmobility

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